

requiring the storage of $O(d)$ cells. Since each cell has $|\Sigma| + 1$ entries, the space complexity is in $O(\min(d^2|\Sigma| \log r, d(n-d)|\Sigma| \log |\Sigma|))$ bits.

The algorithm requires a sum of $|\Sigma|$ numbers and $|\Sigma|$ copies per cell. Since all the numbers have $O(\min(d \log r, (n-d) \log |\Sigma|))$ bits and the table has nd cells, the time complexity is in $O(nd|\Sigma| \min(d \log r, (n-d) \log |\Sigma|))$. \square

V. FUTURE WORK

For future work, it would be interesting to find closed-form expressions and an efficient algorithm to calculate the number of subsequences of a string when deletions and insertions of symbols are allowed.

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On Optimal Quantization Rules for Some Problems in Sequential Decentralized Detection

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Abstract—We consider the design of systems for sequential decentralized detection, a problem that entails several interdependent choices: the choice of a stopping rule (specifying the sample size), a global decision function (a choice between two competing hypotheses), and a set of quantization rules (the local decisions on the basis of which the global decision is made). This correspondence addresses an open problem of whether in the Bayesian formulation of sequential decentralized detection, optimal local decision functions can be found within the class of stationary rules. We develop an asymptotic approximation to the optimal cost of stationary quantization rules and exploit this approximation to show that stationary quantizers are not optimal in a broad class of settings. We also consider the class of blockwise-stationary quantizers, and show that asymptotically optimal quantizers are likelihood-based threshold rules.

Index Terms—Decentralized detection, decision-making under constraints, experimental design, hypothesis testing, quantizer design, sequential detection.

I. INTRODUCTION

Detection is a classical discrimination or hypothesis-testing problem, in which observations $\{X_1, X_2, \dots\}$ are assumed to be drawn independent and identically distributed (i.i.d.) from the (multivariate) conditional distribution $\mathbb{P}(\cdot | H)$ and the goal is to infer the value of the random variable H , which takes values in $\{0, 1\}$. In a typical engineering application, the case $\{H = 1\}$ represents the presence of some target to be detected, whereas $\{H = 0\}$ represents its absence. Placing this problem in a communication-theoretic context, a *decentralized detection* problem is a hypothesis-testing problem in which the decision maker is not given access to the raw data points X_n , but instead must infer H based only on the output of a set of *quantization rules* or *local decision functions*, say $\{U_n = \phi_n(X_n)\}$, which map the raw data to quantized values. This basic problem of decentralized detection has been studied extensively for several decades [17], [19], [6]; see the overview papers [20], [23], [3], [5], and references therein, for more background. Of interest in this correspondence is the extension to an-online setting: more specifically, the *sequential decentralized detection* problem [19], [21], [12] involves a data sequence $\{X_1, X_2, \dots\}$, and a corresponding sequence of summary statistics $\{U_1, U_2, \dots\}$, determined by a sequence of local decision rules $\{\phi_1, \phi_2, \dots\}$. The goal is to design both the local decision functions and to specify a global decision rule so as to predict H in a manner that optimally trades off accuracy and delay. In short, the sequential decentralized detection problem is the communication-constrained extension of classical formulation of sequential

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centralized decision-making problems (see, e.g., [8], [15], [10]) to the decentralized setting.

In setting up a general framework for studying sequential decentralized problems, Veeravalli *et al.* [22] defined five problems, denoted “Case A” through “Case E,” distinguished from one another by the amount of information available to the local sensors. In applications such as power-constrained sensor networks, one cannot assume that the decision maker and sensors can communicate over a high-bandwidth channel, nor that the sensors have unbounded memory. Most suited to this perspective—and the focus of this correspondence—is Case A, in which the local decisions are of the simplified form $\phi_n(X_n)$; i.e., neither local memory nor feedback are assumed to be available. Noting that Case A is not amenable to dynamic programming and hence presumably intractable, Veeravalli *et al.* [22] suggested restricting the analysis to the class of *stationary* local decision functions; i.e., local decision functions ϕ_n that are independent of n . They conjectured that stationary decision functions might actually be optimal in the setting of Case A (given the intuitive symmetry and high degree of independence of the problem in this case), even though it is not possible to verify this optimality via dynamic programming (DP) arguments. This conjecture has remained open since it was first posed by Veeravalli *et al.* [22], [21].

The main contribution of this correspondence is to resolve this question by showing that stationary decision functions are, in fact, *not* optimal for decentralized problems of type A. Our argument is based on an asymptotic characterization of the optimal Bayesian risk as the cost per sample goes to zero. In this asymptotic regime, the optimal cost can be expressed as a simple function of priors and Kullback–Leibler (KL) divergences. This characterization allows us to construct counterexamples to the stationarity conjecture, both in an exact and an asymptotic setting. In the latter setting, we present a class of problems in which there always exists a range of prior probabilities for which stationary strategies, either deterministic or randomized, are suboptimal. We note in passing that an intuition for the source of the suboptimality is easily provided—it is due to the asymmetry of the KL divergence.

It is well known that optimal quantizers when unrestricted are necessarily likelihood-based threshold rules [19]. Our counterexamples and analysis imply that optimal thresholds are not generally stationary (i.e., the threshold may differ from sample to sample). We also provide a partial converse to this result: specifically, if we restrict ourselves to stationary (or blockwise stationary) quantizer designs, then there exists an optimal design that is a deterministic threshold rule based on the likelihood ratio. We prove this result by establishing a quasi-convexity result for the asymptotically optimal cost function.

It is worth highlighting several limitations in our results. For the suboptimality of stationary quantizers, our analysis is applicable only to finite classes of deterministic quantizers and their convex hull of randomized quantizers, and under the assumption that the likelihood ratio of the two hypotheses are bounded from both above and below. Such assumptions certainly hold for arbitrary discrete distributions with finite support. It remains an open problem to consider more general classes of distributions. For the likelihood-ratio characterization result, our proof works only for the (possibly infinite) classes of deterministic quantizers with arbitrary output alphabets, as well as for the class of randomized quantizers with binary outputs. We conjecture that the same result holds more generally for randomized quantizers with arbitrary output alphabets.

The remainder of this correspondence is organized as follows. We begin in Section II with background on the Bayesian formulation of sequential detection problems, and Wald’s approximation. Section III provides a simple asymptotic approximation of the optimal cost that underlies our main analysis in Section IV. In Section V, we establish the existence of optimal decision rules that are likelihood-based threshold rules, under the restriction to blockwise stationarity. We conclude with a discussion in Section VI.

II. BACKGROUND

This section provides background on the Bayesian formulation of sequential (centralized) detection problems. Of particular use in our subsequent analysis is Wald’s approximation of the cost of optimal sequential test.

Let \mathbb{P}_0 and \mathbb{P}_1 represent the distribution of X , when conditioned on $\{H = 0\}$ and $\{H = 1\}$, respectively. Assume that \mathbb{P}_0 and \mathbb{P}_1 are absolutely continuous with respect to one another. We use $f^0(x)$ and $f^1(x)$ to denote the respective density functions with respect to some dominating measure (e.g., Lebesgue for continuous variables, or counting measure for discrete-valued variables).

Our focus is the Bayesian formulation of the sequential detection problem [15], [21]; accordingly, we let $\pi^1 = \mathbb{P}(H = 1)$ and $\pi^0 = \mathbb{P}(H = 0)$ denote the prior probabilities of the two hypotheses. Let X_1, X_2, \dots be a sequence of conditionally i.i.d. realizations of X . A sequential decision rule consists of a *stopping time* N defined with respect to the sigma field $\sigma(X_1, \dots, X_N)$, and a decision function γ measurable with respect to $\sigma(X_1, \dots, X_N)$. The cost function is the expectation of a weighted sum of the sample size N and the probability of incorrect decision—namely

$$J(N, \gamma) := \mathbb{E}\{cN + \mathbb{I}[\gamma(X_1, \dots, X_N) \neq H]\} \quad (1)$$

where $c > 0$ is the incremental cost of each sample. The overall goal is to choose the pair (N, γ) so as to minimize the expected loss (1).

It is well known that the optimal solution of the sequential decision problem can be characterized recursively using DP arguments [1], [25], [15], [2]. Although useful in classical (centralized) sequential detection, the DP approach is not always straightforward to apply to *decentralized* versions of sequential detection [21]. In the remainder of this section, we describe an asymptotic approximation of the optimal sequential cost, originally due to Wald (cf. [16]), valid as $c \rightarrow 0$. To sketch out Wald’s approximation, we begin by noting the optimal stopping rule for the cost function (1) takes the form

$$N = \inf \left\{ n \geq 1 \mid L_n(X_1, \dots, X_n) := \sum_{i=1}^n \log \frac{f^1(X_i)}{f^0(X_i)} \notin (a, b) \right\} \quad (2)$$

for some real numbers $a < b$. Given this stopping rule, the optimal decision function has the form

$$\gamma(L_N) = \begin{cases} 1, & \text{if } L_N \geq b \\ 0, & \text{if } L_N \leq a. \end{cases} \quad (3)$$

Consider the two types of error

$$\begin{aligned} \alpha &= \mathbb{P}_0(\gamma(L_N) \neq H) = \mathbb{P}_0(L_N \geq b) \\ \beta &= \mathbb{P}_1(\gamma(L_N) \neq H) = \mathbb{P}_1(L_N \leq a). \end{aligned}$$

As $c \rightarrow 0$, it can be shown that the optimal choice of a and b satisfies $a \rightarrow -\infty$, $b \rightarrow \infty$, and the corresponding α, β satisfy $\alpha + \beta \rightarrow 0$. Ignoring the overshoot of L_N upon the optimal stopping time N (i.e., instead assuming L_N attains precisely the value a or b) we can express $a, b, \mathbb{E}N$, and the cost function J in terms of α and β as follows [24]:

$$a \approx \log \frac{\beta}{1 - \alpha} \quad \text{and} \quad b \approx \log \frac{1 - \beta}{\alpha} \quad (4)$$

$$\mathbb{E}_0[L_N] \approx (1 - \alpha)a + \alpha b \quad (5)$$

$$\mathbb{E}_1[L_N] \approx (1 - \beta)b + \beta a. \quad (6)$$

Now define the KL divergences

$$D^1 = \mathbb{E}_1 \left[\log \frac{f^1(X_1)}{f^0(X_1)} \right] = D(f^1 \parallel f^0)$$

$$D^0 = -\mathbb{E}_0 \left[\log \frac{f^1(X_1)}{f^0(X_1)} \right] = D(f^0 \parallel f^1).$$

With a slight abuse of notation, we shall also use $D(\alpha, \beta)$ to denote a function in $[0, 1]^2 \rightarrow \mathbb{R}$ such that

$$D(\alpha, \beta) := \alpha \log \frac{\alpha}{\beta} + (1 - \alpha) \log \frac{1 - \alpha}{1 - \beta}.$$

With the above approximations, the cost function J of the decision rule based on envelopes a and b can be written as

$$J = \pi^1 \mathbb{E}_1(cN + \mathbb{1}[L_N \leq a]) + \pi^0 \mathbb{E}_0(cN + \mathbb{1}[L_N \geq b]) \quad (7)$$

$$= c\pi^1 \frac{\mathbb{E}_1 L_N}{D^1} + c\pi^0 \frac{\mathbb{E}_0 L_N}{-D^0} + \pi^0 \alpha + \pi^1 \beta \quad (8)$$

$$\approx c\pi^0 \frac{D(\alpha, 1 - \beta)}{D^0} + c\pi^1 \frac{D(1 - \beta, \alpha)}{D^1} + \pi^0 \alpha + \pi^1 \beta \quad (9)$$

where the last approximation follows from Wald's equation [24]. Let $\tilde{J}(\alpha, \beta)$ denote approximation (9) of J .

Let J^* denote the cost of an optimal sequential test

$$J^* = \inf_{a, b} J, \quad (10)$$

A useful result due Chernoff [7] states that under certain assumption (to be elaborated on in the next section), J^* has the following form:

$$J^* \approx \left(\frac{\pi^0}{D^0} + \frac{\pi^1}{D^1} \right) c \log c^{-1} (1 + o(1)). \quad (11)$$

III. CHARACTERIZATION OF OPTIMAL STATIONARY QUANTIZERS

Turning now to the decentralized setting, the primary challenge lies in the design of the quantization rules ϕ_n applied to data X_n . When X_n is univariate, a deterministic quantization rule ϕ_n is a function that maps \mathcal{X} to the discrete space $\mathcal{U} = \{0, \dots, K - 1\}$ for some natural number K . For multivariate X_n with d dimensions arising in the multiple-sensor setting, a deterministic quantizer ϕ_n is defined as a mapping from the d -dimensional product space \mathcal{X} to $\mathcal{U} = \{0, \dots, K - 1\}^d$. In the decentralized problem defined as Case A by Veeravalli *et al.* [22], the function ϕ_n is composed of d separate quantizer functions, one each for each dimension. A randomized quantizer ϕ_n is obtained by placing a distribution over the space of deterministic quantizers.

Any fixed set of quantization rules ϕ_n yields a sequence of compressed data $U_n = \phi_n(X_n)$, to which the classical theory can be applied. We are thus interested in choosing quantization rules ϕ_1, ϕ_2, \dots so that the error resulting from applying the optimal sequential test to the sequence of statistics U_1, U_2, \dots is minimized over some space Φ of quantization rules. For a given quantizer ϕ_n we use

$$f_{\phi_n}^i(u) := \mathbb{P}_i(\phi_n(X_n) = u), \quad \text{for } i = 0, 1$$

to denote the distributions of the compressed data, conditioned on the hypothesis. In general, when randomized quantizers are allowed, the vector $(f_{\phi_n}^0(\cdot), f_{\phi_n}^1(\cdot))$ ranges over a convex set, denoted $\text{conv}\Phi$, whose extreme points correspond to deterministic quantizers based on likelihood-ratio threshold rules [18].

We say that a quantizer design is *stationary* if the rule ϕ_n is independent of n ; in this case, we simplify the notation to f_ϕ^1 and f_ϕ^0 . In addition, we define the KL divergences $D_\phi^1 := D(f_\phi^1 \| f_\phi^0)$ and $D_\phi^0 := D(f_\phi^0 \| f_\phi^1)$. Moreover, let J_ϕ and J_ϕ^* denote the analogues of the functions J in (7) and J^* in (10), respectively, defined using D_ϕ^i , for $i = 0, 1$. In this scenario, the sequence of compressed data U_1, \dots, U_n, \dots are drawn i.i.d. from either f_ϕ^0 or f_ϕ^1 . Thus, we can use the approximation (11) to characterize the asymptotically optimal stationary quantizer design. This is stated formally in the lemma to follow.

We begin by stating the assumptions underlying the lemma. For a given class of quantizers Φ , we assume that the KL divergences are uniformly bounded away from zero

$$D(f_\phi^1 \| f_\phi^0) > 0, \quad D(f_\phi^0 \| f_\phi^1) > 0, \quad \text{for all } \phi \in \Phi \quad (12)$$

and, moreover, that the variance of the log-likelihood ratios are bounded

$$\sup_{\phi \in \Phi} \max \left\{ \text{Var}_{f_\phi^1} \log(f_\phi^1 / f_\phi^0), \text{Var}_{f_\phi^0} \log(f_\phi^0 / f_\phi^1) \right\} < \infty. \quad (13)$$

Lemma 1:

(a) Under assumptions (12) and (13), the optimal stationary cost takes the form

$$J_\phi^* = \left(\frac{\pi^0}{D_\phi^0} + \frac{\pi^1}{D_\phi^1} \right) c \log c^{-1} (1 + r_\phi) \quad (14)$$

where $|r_\phi| = o(1)$ as $c \rightarrow 0$.

(b) If

$$\sup_{\phi \in \Phi} \max \{ \log(f_\phi^1 / f_\phi^0), \log(f_\phi^0 / f_\phi^1) \} < M$$

for some constant M , then (14) holds with $\sup_{\phi \in \Phi} |r_\phi| = o(1)$ as $c \rightarrow 0$.

Proof:

(a) This part is immediate from a combination of Theorems 1 and 2 of Chernoff [7].

(b) We begin by bounding the error in the approximation (9). By definition of the stopping time N , we have either i) $b \leq L_N \leq b + M$ or ii) $a - M \leq L_N \leq a$. By standard arguments due to Wald [24], it is simple to obtain $e^b \alpha \leq 1 - \beta \leq e^{b+M} \alpha$, or equivalently, $b \leq b(\alpha, \beta) = \log \frac{1 - \beta}{\alpha} \leq b + M$. Similar reasoning for case ii) yields $a - M \leq a(\alpha, \beta) = \log \frac{\beta}{1 - \alpha} \leq a$. Now, note that

$$\mathbb{E}_0 L_N = \alpha \mathbb{E}_0[L_N | L_N \geq b] + (1 - \alpha) \mathbb{E}_0[L_N | L_N \leq a].$$

Conditioning on the event $L_N \in [b, b + M]$, we have $|L_N - b(\alpha, \beta)| \leq M$. Similarly, conditioning on the event $L_N \in [a - M, a]$, we have $|L_N - a(\alpha, \beta)| \leq M$. This yields $|\mathbb{E}_0 L_N - (-D(\alpha, 1 - \beta))| \leq M$. Similar reasoning yields $|\mathbb{E}_1 L_N - D(1 - \beta, \alpha)| \leq M$. Let $\tilde{J}_\phi(a, b)$ denote the approximation (9) of J_ϕ . We obtain

$$|J_\phi - \tilde{J}_\phi(\alpha, \beta)| \leq 2cM.$$

Note that the approximation error bound is independent of ϕ . Thus, it suffices to establish the asymptotic behavior (14) for the quantity $\inf_{\alpha, \beta} \tilde{J}_\phi(\alpha, \beta)$, where the infimum is taken over pairs of realizable error probabilities (α, β) . Moreover, we only need to consider the asymptotic regime $\alpha + \beta \rightarrow 0$, since the error probabilities α and β vanish as $c \rightarrow 0$. It is simple to see that

$$D(1 - \beta, \alpha) = \log(1/\alpha)(1 + o(1))$$

and

$$D(1 - \alpha, \beta) = \log(1/\beta)(1 + o(1)).$$

Hence, $\inf_{\alpha, \beta} \tilde{J}_\phi(\alpha, \beta)$ can be expressed as

$$\inf_{\alpha, \beta} \left\{ \pi^0 \alpha + \pi^1 \beta + c\pi^0 \frac{\log(1/\beta)}{D_\phi^0} + c\pi^1 \frac{\log(1/\alpha)}{D_\phi^1} \right\} (1 + o(1)). \quad (15)$$

This infimum, taken over all positive (α, β) , is achieved at

$$\alpha^* = \frac{c\pi^1}{D_\phi^1 \pi^0} \quad \text{and} \quad \beta^* = \frac{c\pi^0}{D_\phi^0 \pi^1}.$$

Plugging the quantities α^* and β^* into (15) yields (14). Note that the asymptotic quantity $o(1)$ in (14) is absolutely bounded by $\alpha^* + \beta^* \rightarrow 0$ uniformly for all quantizer ϕ , because D_ϕ^1 and D_ϕ^0 are uniformly bounded away from zero due to the lemma's assumption.

It remains to show that error probabilities (α^*, β^*) can be approximately realized by using a sufficiently large threshold $b > 0$ and small threshold $a < 0$ while incurring an approximation cost of order $O(c)$ uniformly for all ϕ . Indeed, let us choose thresholds a' and b' such that $e^{-(b'+M)}/2 \leq \alpha^* \leq e^{-b'}$, and $e^{a'-M}/2 \leq \beta^* \leq e^{a'}$. Let α' and β' be the corresponding errors associated with these two thresholds. As before, we also have $\alpha' \in (e^{-(b'+M)}/2, e^{-b'})$ and $\beta' \in (e^{a'-M}/2, e^{a'})$. Clearly, $|\alpha^* - \alpha'| \leq e^{-b'}(1 - e^{-M}/2) = O(\alpha^*) = O(c)$. Similarly, $|\beta^* - \beta'| = O(c)$. By the mean value theorem

$$\begin{aligned} |\log(1/\alpha^*) - \log(1/\alpha')| &\leq |\alpha^* - \alpha'|e^{b'+M} \\ &\leq 2e^M(1 - e^{-M}/2) = O(1). \end{aligned} \quad (16)$$

Similarly, $\log(1/\beta^*) - \log(1/\beta') = O(1)$. Hence, the approximation of (α^*, β^*) by the realizable (α', β') incurs a cost at most $O(c)$. Furthermore, the constant in the asymptotic bound $O(c)$ is independent of quantizer $\phi \in \Phi$. \square

Remarks: i) If Φ is a finite class of quantizers, or a convex hull of a finite class of quantizers, the assumption in Part b of Lemma 1 holds. It also holds in the case of discrete distributions and continuous distributions with bounded support. However, it would be interesting to relax this assumption so as to cover distributions with unbounded support. ii) The preceding approximation of the optimal cost essentially ignores the overshoot of the likelihood ratio L_N . While it is possible to analyze this overshoot to obtain a finer approximation (cf. [11], [16], [10], [14]), we see that this is not needed for our purpose. Lemma 1 shows that given a fixed prior (π^0, π^1) , among all stationary quantizer designs in Φ , ϕ is optimal for sufficiently small c if and only if ϕ minimizes what we shall call the *sequential cost coefficient*

$$G_\phi := \frac{\pi^0}{D_\phi^0} + \frac{\pi^1}{D_\phi^1}.$$

iii) As a consequence of Lemma 7 to be proved in the sequel, if we consider the class Φ of all binary randomized quantizers, then sequential cost coefficient G_ϕ is a quasi-concave function with respect to $(f_\phi^0(\cdot), f_\phi^1(\cdot))$. (A function F is quasi-concave if and only if for any η , the level set $\{F(x) \geq \eta\}$ is a convex set; see Boyd and Vandenberghe [4] for further background). The minimum of a quasi-concave function lies in the set of extreme points in its domain. For the set $\text{conv}\Phi$, these extreme points can be realized by deterministic quantizers based on likelihood ratios [20]. Consequently, we conclude that for quantizers with binary outputs, the optimal cost is not decreased by considering randomized quantizers. We conjecture that this statement also holds beyond the binary case.

Section V is devoted to a more detailed study of asymptotically optimal stationary quantizers. In the meantime, we turn to the question of whether stationary quantizers are optimal in either finite-sample or asymptotic settings.

IV. SUBOPTIMALITY OF STATIONARY DESIGNS

It was shown by Tsitsiklis [19] that optimal quantizers ϕ_n take the form of threshold rules based on the likelihood ratio $f^1(X_n)/f^0(X_n)$. Veeravalli *et al.* [22], [21] asked whether these rules can always be taken to be stationary, a conjecture that has remained open. In this section, we resolve this question with a negative answer in both the finite-sample and asymptotic settings.

A. Suboptimality in Exact Setting

We begin by providing a numerical counterexample for which stationary designs are suboptimal. Consider a problem in which $X \in \mathcal{X} = \{1, 2, 3\}$ and the conditional distributions take the form

$$f^0(x) = \left[\frac{8}{10} \quad \frac{1999}{10000} \quad \frac{1}{10000} \right] \text{ and } f^1(x) = \left[\frac{1}{3} \quad \frac{1}{3} \quad \frac{1}{3} \right].$$

Suppose that the prior probabilities are $\pi^1 = \frac{8}{100}$ and $\pi^0 = \frac{92}{100}$, and that the cost for each sample is $c = \frac{1}{100}$.

If we restrict to binary quantizers (i.e., $\mathcal{U} = \{0, 1\}$), by the symmetric roles of the output alphabets, there are only three possible deterministic quantizers.

- 1) Design A: $\phi_A(X_n) = 0 \iff X_n = 1$. As a result, the corresponding distribution for U_n is specified by $f_{\phi_A}^0(u_n) = [\frac{4}{5} \quad \frac{1}{5}]$ and $f_{\phi_A}^1(u) = [\frac{1}{3} \quad \frac{2}{3}]$.
- 2) Design B: $\phi_B(X_n) = 0 \iff X_n \in \{1, 2\}$. The corresponding distribution for U_n is given by $f_{\phi_B}^0(u) = [\frac{9999}{10000} \quad \frac{1}{10000}]$ and $f_{\phi_B}^1(u) = [\frac{2}{3} \quad \frac{1}{3}]$.
- 3) Design C: $\phi_C(X_n) = 0 \iff X_n \in \{1, 3\}$. The corresponding distribution for U_n is specified by $f_{\phi_C}^0 \sim [\frac{8001}{10000} \quad \frac{1999}{10000}]$ and $f_{\phi_C}^1(u) = [\frac{2}{3} \quad \frac{1}{3}]$.

Now consider the three stationary strategies, each of which uses only one fixed design, A, B, or C. For any given stationary quantization rule ϕ , we have a classical centralized sequential problem, for which the optimal cost (achieved by a sequential probability ratio test) can be computed using a dynamic-programming procedure [25], [1]. Accordingly, for each stationary strategy, we compute the optimal cost function J for 10^6 points on the p -axis by performing 300 updates of Bellman's equation (cf. [2]). In all cases, the difference in cost between the 299th and 300th updates is less than 10^{-6} . Let J_A, J_B and J_C denote the optimal cost function for sequential tests using all A's, all B's, and all C's, respectively. When evaluated at $\pi^1 = 0.08$, these computations yield $J_A = 0.0567, J_B = 0.0532$, and $J_C = 0.08$.

Finally, we consider a nonstationary rule obtained by applying design A for only the first sample, and applying design B for the remaining samples. Again, using Bellman's equation, we find that the cost for this design is

$$\begin{aligned} J_* &= \min\{\min\{\pi^1, 1 - \pi^1\} \\ &+ J_B(P(H = 1|u_1 = 0))P(u_1 = 0) \\ &+ J_B(P(H = 1|u_1 = 1))P(u_1 = 1)\} = 0.052767 \end{aligned}$$

which is better than any of the stationary strategies.

In this particular example, the cost J^* of the nonstationary quantizer yields a slim improvement (0.0004) over the best stationary rule J_B . This slim margin is due in part to the choice of a small per-sample cost $c = 0.01$; however, larger values of c do not yield a counterexample when using the particular distributions specified above. A more significant factor is that our nonstationary rule differs from the optimal stationary rule B only in its treatment of the first sample. This fact suggests that one might achieve better cost by alternating between using design A and design B on the odd and even samples, respectively. Our analysis of the asymptotic setting in the next section confirms this intuition.

B. Asymptotic Suboptimality for Both Deterministic and Randomized Quantizers

We now prove that in a broad class of examples, there is a range of prior probabilities for which stationary quantizer designs are suboptimal. Our result stems from the following observation: Lemma 1 implies that in order to achieve a small cost we need to choose a quantizer ϕ for which the KL divergences $D_\phi^0 := D(f_\phi^0||f_\phi^1)$ and $D_\phi^1 := D(f_\phi^1||f_\phi^0)$ are both as large as possible. Due to the asymmetry of the KL divergence, however, these maxima are not necessarily achieved by

a single quantizer ϕ . This suggests that one could improve upon stationary designs by applying different quantizers to different samples, as the following lemma shows.

Lemma 2: Let ϕ_1 and ϕ_2 be any two quantizers. If the following inequalities hold:

$$D_{\phi_1}^0 < D_{\phi_2}^0 \quad \text{and} \quad D_{\phi_1}^1 > D_{\phi_2}^1 \quad (17)$$

then there exists a nonempty interval $(U, V) \subseteq (0, +\infty)$ such that as $c \rightarrow 0$

$$\begin{aligned} J_{\phi_1}^* &\leq J_{\phi_1, \phi_2}^* \leq J_{\phi_2}^*, & \text{if } \frac{\pi^0}{\pi^1} \leq U \\ J_{\phi_1, \phi_2}^* &< \min\{J_{\phi_1}^*, J_{\phi_2}^*\} - \Theta(c \log c^{-1}), & \text{if } \frac{\pi^0}{\pi^1} \in (U, V) \\ J_{\phi_1}^* &\geq J_{\phi_1, \phi_2}^* \geq J_{\phi_2}^*, & \text{if } \frac{\pi^0}{\pi^1} \geq V \end{aligned}$$

where J_{ϕ_1, ϕ_2}^* denotes the optimal cost of a sequential test that alternates between using ϕ_1 and ϕ_2 on odd and even samples respectively.

Proof: According to Lemma 1, for $i = 0, 1$

$$J_{\phi_i}^* = \left(\frac{\pi^0}{D_{\phi_i}^0} + \frac{\pi^1}{D_{\phi_i}^1} \right) c \log c^{-1} (1 + o(1)). \quad (18)$$

Now consider the sequential test that applies quantizers ϕ_1 and ϕ_2 alternately to odd and even samples. Furthermore, let this test consider two samples at a time. Let $f_{\phi_1 \phi_2}^0$ and $f_{\phi_1 \phi_2}^1$ denote the induced conditional probability distributions, jointly on the odd–even pairs of quantized variables. From the additivity of the KL divergence and assumption (17), there holds

$$D(f_{\phi_1 \phi_2}^0 \| f_{\phi_1 \phi_2}^1) = D_{\phi_1}^0 + D_{\phi_2}^0 > 2D_{\phi_1}^0 \quad (19a)$$

$$D(f_{\phi_1 \phi_2}^1 \| f_{\phi_1 \phi_2}^0) = D_{\phi_1}^1 + D_{\phi_2}^1 < 2D_{\phi_1}^1. \quad (19b)$$

Clearly, the cost of the proposed sequential test is an upper bound for J_{ϕ_1, ϕ_2}^* . Furthermore, the gap between this upper bound and the true optimal cost is no more than $O(c)$. Hence, as in the proof of Lemma 1, as $c \rightarrow 0$, the optimal cost J_{ϕ_1, ϕ_2}^* can be written as

$$\left(\frac{2\pi^0}{D_{\phi_1}^0 + D_{\phi_2}^0} + \frac{2\pi^1}{D_{\phi_1}^1 + D_{\phi_2}^1} \right) c \log c^{-1} (1 + o(1)). \quad (20)$$

From (18) and (20), simple calculations yield the claim with

$$\begin{aligned} U &= \frac{D_{\phi_1}^0 (D_{\phi_1}^1 - D_{\phi_2}^1) (D_{\phi_1}^0 + D_{\phi_2}^0)}{D_{\phi_1}^1 (D_{\phi_1}^1 + D_{\phi_2}^1) (D_{\phi_2}^0 - D_{\phi_1}^0)} \\ &< V = \frac{D_{\phi_2}^0 (D_{\phi_1}^1 - D_{\phi_2}^1) (D_{\phi_1}^0 + D_{\phi_2}^0)}{D_{\phi_2}^1 (D_{\phi_1}^1 + D_{\phi_2}^1) (D_{\phi_2}^0 - D_{\phi_1}^0)}. \end{aligned} \quad (21)$$

□

Example: Let us return to the example provided in the previous subsection. Note that the two quantizers ϕ_A and ϕ_B satisfy assumption (17), since $D(f_{\phi_B}^0 \| f_{\phi_B}^1) = 0.4045 < D(f_{\phi_A}^0 \| f_{\phi_A}^1) = 0.45$ and $D(f_{\phi_B}^1 \| f_{\phi_B}^0) = 2.4337 > D(f_{\phi_A}^1 \| f_{\phi_A}^0) = 0.5108$. Furthermore, both quantizers dominate ϕ_C in terms of KL divergences: $D(f_{\phi_C}^0 \| f_{\phi_C}^1) = 0.0438$, $D(f_{\phi_C}^1 \| f_{\phi_C}^0) = 0.0488$. As a result, there exist a range of priors for which a sequential test using stationary quantizer design (either ϕ_A , ϕ_B or ϕ_C for all samples) is not optimal.

Theorem 3:

(a) Suppose that Φ is a finite collection of quantizers, and that there is no single quantizer ϕ that dominates all other quantizers in Φ in the

sense that

$$D_{\phi}^0 \geq D_{\phi'}^0 \quad \text{and} \quad D_{\phi}^1 \geq D_{\phi'}^1, \quad \text{for all } \phi' \in \Phi. \quad (22)$$

Then there exists a nonempty range of prior probabilities for which no stationary design based on a quantizer in Φ is optimal.

(b) For any nondeterministic ϕ in the randomized class $\text{conv} \Phi$, there exists a nonstationary quantizer design that has strictly smaller sequential cost coefficient than that of a stationary design based on ϕ for any choice of prior probabilities.

Proof:

(a) Since there are a finite number of quantizers in Φ and no quantizer dominates all others, the interval $(0, \infty)$ is divided into at least two adjacent nonempty intervals, each of which corresponds to a range of prior probability ratios π^0/π^1 for which a quantizer is strictly optimal (asymptotically) among all stationary designs. Let them be (δ_1, δ) and (δ, δ_2) , for two quantizers, namely, ϕ_1 and ϕ_2 . In particular, δ is the value for π^0/π^1 for which the sequential cost coefficients are equal—viz. $G_{\phi_1} = G_{\phi_2}$ —which happens only if assumption (17) holds. Some calculations verify that

$$\delta = \frac{D_{\phi_1}^0 D_{\phi_2}^0 (D_{\phi_2}^1 - D_{\phi_1}^1)}{D_{\phi_1}^1 D_{\phi_2}^1 (D_{\phi_1}^0 - D_{\phi_2}^0)}. \quad (23)$$

By Lemma 2, a nonstationary design obtained by alternating between ϕ_1 and ϕ_2 has smaller sequential cost than both ϕ_1 and ϕ_2 for $\pi^0/\pi^1 \in (U, V)$, where U and V are given in (21). Since it can be verified that δ as defined (23) belongs to the interval (U, V) , we conclude that for $\pi^0/\pi^1 \in (U, V) \cap (\delta_1, \delta_2)$, this nonstationary design has smaller cost than any stationary design using $\phi \in \Phi$.

(b) Let $\phi \in \text{conv} \Phi$ be a randomized quantizer (i.e., at each step choose with nonzero probabilities w_1, \dots, w_k from quantizers $\phi_1, \dots, \phi_k \in \Phi$, respectively, where $\sum_{i=1}^k w_i = 1$). Clearly, the density induced by ϕ satisfies $f_{\phi}^0 = \sum_{i=1}^k w_i f_{\phi_i}^0$ and $f_{\phi}^1 = \sum_{i=1}^k w_i f_{\phi_i}^1$. Due to strict convexity of the KL divergence functional with respect jointly to the two density arguments [9], by Jensen's inequality we have $D_{\phi}^0 < \sum_{i=1}^k w_i D_{\phi_i}^0$ and $D_{\phi}^1 < \sum_{i=1}^k w_i D_{\phi_i}^1$. Since $D_{\phi_i}^0$ and $D_{\phi_i}^1$ are bounded from above uniformly for all $\phi_i \in \Phi$, it is possible to approximate (w_1, \dots, w_k) by rational numbers of the form $(q_1/N, q_2/N, \dots, q_k/N)$ for some natural numbers q_1, \dots, q_k and N satisfying $\sum_{i=1}^k q_i = N$ such that

$$\begin{aligned} D_{\phi}^0 &< \sum_{i=1}^k q_i D_{\phi_i}^0 / N \\ D_{\phi}^1 &< \sum_{i=1}^k q_i D_{\phi_i}^1 / N. \end{aligned}$$

Now consider the nonstationary quantizer that applies ϕ_1 for q_1 steps, then ϕ_2 for q_2 steps, and so on, up to ϕ_k for q_k steps, yielding a total of N steps, and then repeats this sequence starting again at step $N + 1$. By construction, this nonstationary quantizer has a smaller cost than that of quantizer ϕ for any choice of prior. □

Remarks: i) It is worth emphasizing the assumption that the class Φ is finite is crucial in Part (a) of the theorem. We do not know if this result can be extended to the case in which Φ is infinite. ii) Part (b) shows that any stationary randomized quantizer is always dominated by some nonstationary one. Actually, a stronger result can be proved at least for binary quantizers (see Corollary 8): for any given choice of prior probability, any stationary randomized quantizer is dominated by a stationary deterministic quantizer. iii) It is interesting to contrast the Bayesian formulation of the problem of quantizer design with the Neyman–Pearson formulation. Our results on the suboptimality of stationary quantizer

design in the Bayesian formulation repose on the asymmetry of the KL divergence, as well as the sensitivity of the optimal quantizers on the prior probability. We note that Mei [12 (see p. 58)] considered the Neyman–Pearson formulation of this problem. In this formulation, it can be shown that for all sequential tests for which the Type 1 and Type 2 errors are bounded by α and β , respectively, then as $\alpha + \beta \rightarrow 0$, the expected stopping time $\mathbb{E}_0 N$ under hypothesis $H = 0$ is asymptotically minimized by applying a stationary quantizer ϕ^* that maximizes $D(f_\phi^0 \| f_\phi^1)$. Similarly, the expected stopping time $\mathbb{E}_1 N$ under hypothesis $H = 1$ is asymptotically minimized by the stationary quantizer ϕ^{**} that maximizes $D(f_\phi^1 \| f_\phi^0)$ [12]. In this context, the example in Subsection IV-A provides a case in which the asymptotically minimal KL divergences ϕ^* and ϕ^{**} are not the same, due to the asymmetry, which suggests that there may not exist a stationary quantizer that simultaneously minimizes both $\mathbb{E}_1 N$ and $\mathbb{E}_0 N$.

C. Asymptotic Suboptimality in Multiple-Sensor Setting

Our analysis thus far has established that with a single sensor per time step ($d = 1$), applying multiple quantizers to different samples can reduce the sequential cost. As pointed out by one of the referees, it is natural to ask whether the same phenomenon persists in the case of multiple sensors ($d > 1$). In this section, we show that the phenomenon does indeed carry over, more specifically by providing an example in which stationary strategies are still suboptimal in comparison to nonstationary ones. The key insight is that we have only a fixed number of dimensions, whereas as $c \rightarrow 0$ we are allowed to take more samples, and each sample can act as an extra dimension, providing more flexibility for nonstationary strategies.

Suppose that the observation vector X_n at time n is d -dimensional, with each component corresponding to a sensor in a typical decentralized setting. Suppose that the observations from each sensor are assumed to be i.i.d. according to the conditional distributions defined in our earlier example (see Section IV-A). Of interest are the optimal deterministic binary quantizer designs for all d sensors. Although there are three possible choices ϕ_A, ϕ_B , and ϕ_C for each sensor, the quantizer ϕ_C is dominated by the other two, so each sensor should choose either ϕ_A and ϕ_B . Suppose that among these sensors, a subset of size k choose ϕ_A and whereas the remaining $d - k$ sensors choose ϕ_B for $0 \leq k \leq d$. We thus have $d + 1$ possible stationary designs to consider. For each k , the sequential cost coefficient corresponding to the associated stationary design takes the form

$$G_k := \frac{\pi^0}{kD_{\phi_A}^0 + (d-k)D_{\phi_B}^0} + \frac{\pi^1}{kD_{\phi_A}^1 + (d-k)D_{\phi_B}^1}. \quad (24)$$

Now consider the following nonstationary design: the first sensor alternates between decision rules ϕ_A and ϕ_B , while the remaining $d - 1$ sensors simply apply the stationary design based on ϕ_B . For this design, the associated sequential cost coefficient is given by

$$G := \frac{2\pi^0}{D_{\phi_A}^0 + (2d-1)D_{\phi_B}^0} + \frac{2\pi^1}{D_{\phi_A}^1 + (2d-1)D_{\phi_B}^1}. \quad (25)$$

Consider the interval (U, V) , where the interval has endpoints

$$\begin{aligned} U &= \frac{D_{\phi_B}^1 - D_{\phi_A}^1}{D_{\phi_A}^0 - D_{\phi_B}^0} \frac{D_{\phi_A}^0 + (2d-1)D_{\phi_B}^0}{D_{\phi_A}^1 + (2d-1)D_{\phi_B}^1} \\ &\times \frac{D_{\phi_B}^0}{D_{\phi_B}^1} < V = \frac{D_{\phi_B}^1 - D_{\phi_A}^1}{D_{\phi_A}^0 - D_{\phi_B}^0} \\ &\times \frac{D_{\phi_A}^0 + (2d-1)D_{\phi_B}^0}{D_{\phi_A}^1 + (2d-1)D_{\phi_B}^1} \frac{D_{\phi_A}^0 + (d-1)D_{\phi_B}^0}{D_{\phi_A}^1 + (d-1)D_{\phi_B}^1}. \end{aligned} \quad (26)$$

Straightforward calculations yield that for any prior likelihood $\pi^0/\pi^1 \in (U, V)$, the minimal cost over stationary designs $\min_{k=0, \dots, d} G_k$ is strictly larger than the sequential cost G of the nonstationary design, previously defined in (25).

V. ON ASYMPTOTICALLY OPTIMAL BLOCKWISE STATIONARY DESIGNS

Despite the possible loss in optimality, it is useful to consider some form of stationarity in order to reduce computational complexity of the optimization and decision process. In this section, we consider the class of *blockwise stationary* designs, meaning that there exists some natural number T such that $\phi_{T+1} = \phi_1, \phi_{T+2} = \phi_2$, and so on. For each T , let C_T denote the class of all blockwise stationary designs with period T . We assume throughout the analysis that each decision rule ϕ_n ($n = 1, \dots, T$) satisfies conditions (12) and (13). Thus, as T increases, we have a hierarchy of increasingly rich quantizer classes that will be seen to yield progressively better approximations to the optimal solution.

For a fixed prior (π^0, π^1) and $T > 0$, let (ϕ_1, \dots, ϕ_T) denote a quantizer design in C_T . As before, the cost J_ϕ^* of an asymptotically optimal sequential test using this quantizer design is of order $c \log c^{-1}$ with the sequential cost coefficient

$$G_\phi = \frac{T\pi^0}{D_{\phi_1}^0 + \dots + D_{\phi_T}^0} + \frac{T\pi^1}{D_{\phi_1}^1 + \dots + D_{\phi_T}^1}. \quad (27)$$

G_ϕ is a function of the vector of probabilities introduced by the quantizer $(f_\phi^0(\cdot), f_\phi^1(\cdot))$. We are interested in the properties of a quantization rule ϕ that minimizes G_ϕ .

It is well known that there exist optimal quantizers—*when unrestricted*—that can be expressed as threshold rules based on the log-likelihood ratio (LLR) [19]. Our counterexamples in the preceding sections imply that the thresholds need not be stationary (i.e., the threshold may differ from sample to sample). In the remainder of this section, we address a partial converse to this issue: specifically, if we restrict ourselves to stationary (or blockwise stationary) quantizer designs, then there exists an optimal design consisting of LLR-based threshold rules.

It turns out that the analysis for the case $T > 1$ can be reduced to an analysis that is closely related to our earlier analysis for $T = 1$. Indeed, consider the sequential cost coefficient for the time step $n = 1$, where the rules for the other time steps are held fixed. From (27) we have

$$G_\phi = \frac{T\pi^0}{D_{\phi_1}^0 + s_0} + \frac{T\pi^1}{D_{\phi_1}^1 + s_1}$$

for nonnegative constants s_0 and s_1 . As we will show, our earlier analysis of the sequential cost coefficient, in which $s_0 = s_1 = 0$, carries through to the case in which these values are nonzero. This allows us to provide (in Theorem 9) a characterization of the optimal blockwise stationary quantizer.

Definition 4: The quantizer design function $\phi : \mathcal{X} \rightarrow \mathcal{U}$ is said to be a *likelihood-ratio threshold rule* if there are thresholds $d_0 = -\infty < d_1 < \dots < d_K = +\infty$, and a permutation (u_1, \dots, u_K) of $(0, 1, \dots, K-1)$ such that for $l = 1, \dots, K$, with \mathbb{P}_0 -probability 1, we have

$$\phi(X) = u_l, \quad \text{if } d_{l-1} \leq f^1(X)/f^0(X) \leq d_l.$$

When $f^1(X)/f^0(X) = d_{l-1}$, set $\phi(X) = u_{l-1}$ or $\phi(X) = u_l$ with \mathbb{P}_0 -probability 1.¹

Previous work on the extremal properties of likelihood-ratio-based quantizers guarantees that the KL divergence is maximized by an LLR-based quantizer [18]. In our case, however, the sequential cost coefficient G_ϕ involves a pair of KL divergences, D_ϕ^0 and D_ϕ^1 , which are related to one another in a nontrivial manner. Hence, establishing asymp-

¹This last requirement of the definition is termed the *canonical* likelihood ratio quantizer by Tsitsiklis [18]. Although one could consider performing additional randomization when there are ties, our later results (in particular, Lemma 7) establish that in this case, randomization will not further decrease the optimal cost J_ϕ^* .

otic optimality of LLR-based rules for this cost function does not follow from existing results, but rather requires further understanding of the interplay between these two KL divergences.

The following lemma concerns certain “unnormalized” variants of the KL divergence. Given vectors $a = (a_0, a_1)$ and $b = (b_0, b_1)$, we define functions \tilde{D}^0 and \tilde{D}^1 mapping from \mathbb{R}_+^2 to the real line as follows:

$$\tilde{D}^0(a, b) := a_0 \log \frac{a_0}{a_1} + b_0 \log \frac{b_0}{b_1} \quad (28a)$$

$$\tilde{D}^1(a, b) := a_1 \log \frac{a_1}{a_0} + b_1 \log \frac{b_1}{b_0}. \quad (28b)$$

These functions are related to the standard (normalized) KL divergence via the relations $\tilde{D}^0(a, 1-a) \equiv D(a_0, a_1)$, and $\tilde{D}^1(a, 1-a) \equiv D(a_1, a_0)$.

Lemma 5: For any positive scalars $a_1, b_1, c_1, a_0, b_0, c_0$ such that $\frac{a_1}{a_0} < \frac{b_1}{b_0} < \frac{c_1}{c_0}$, at least one of the two following conditions must hold:

$$(i) \quad \tilde{D}^0(a, b+c) > \tilde{D}^0(b, c+a) \quad \text{and} \quad (29a)$$

$$\tilde{D}^1(a, b+c) > \tilde{D}^1(b, c+a) \quad (29b)$$

$$(ii) \quad \tilde{D}^0(c, a+b) > \tilde{D}^0(b, c+a) \quad \text{and} \quad (29c)$$

$$\tilde{D}^1(c, a+b) > \tilde{D}^1(b, c+a). \quad (29d)$$

This lemma implies that under certain conditions on the ordering of the probability ratios, one can increase *both* KL divergences by re-quantizing. This insight is used in the following lemma to establish that the optimal quantizer ϕ behaves almost like a likelihood-ratio rule. To state the result, recall that the *essential supremum* is the infimum of the set of all η such that $f(x) \leq \eta$ for \mathbb{P}_0 -almost all x in the domain, for a measurable function f .

Lemma 6: If ϕ is an asymptotically optimal quantizer, then for all pairs $(u_1, u_2) \in \mathcal{U}$, $u_1 \neq u_2$, there holds

$$\frac{f^1(u_1)}{f^0(u_1)} \notin \left(\text{ess inf}_{x:\phi(x)=u_2} \frac{f^1(x)}{f^0(x)}, \text{ess sup}_{x:\phi(x)=u_2} \frac{f^1(x)}{f^0(x)} \right).$$

Note that a likelihood-ratio rule guarantees something stronger: For \mathbb{P}_0 -almost all x such that $\phi(x) = u_1$, $f^1(x)/f^0(x)$ takes a value either to the left or to the right, but not to both sides, of the interval specified above.

Lemma 7 stated below essentially guarantees quasi-concavity of G_ϕ for the case of binary quantizers. To state the result, let $F: [0, 1]^2 \rightarrow \mathbb{R}$ be given by

$$F(a_0, a_1) = \frac{c_0}{D(a_0, a_1) + d_0} + \frac{c_1}{D(a_1, a_0) + d_1}. \quad (30)$$

Lemma 7: For any nonnegative constants c_0, c_1, d_0, d_1 , the function F defined in (30) is quasi-concave.

We provide a proof of this result in the Appendix. An immediate consequence of Lemma 7 is that LLR-based quantizers exist for the class of randomized quantizers with binary outputs.

Corollary 8: Restricting to the class of (blockwise) stationary binary quantizers, there exists an asymptotically optimal quantizer ϕ that is a (deterministic) likelihood-ratio threshold rule.

Proof: Let ϕ is a (randomized) binary quantizer. The sequential cost coefficient can be written as $G_\phi = F(f_\phi^0(0), f_\phi^1(0))$. The set of $\{(f_\phi^0(0), f_\phi^1(0))\}$ for all ϕ is a convex set whose extreme points can be realized by deterministic likelihood ratio threshold rules (see Proposition 3.2 of [18]). Since the minimum of a quasi-concave function must

lie at one such extreme point [4], the corollary is immediate as a consequence of Lemma 7. \square

It turns out that the same statement can also be proved for deterministic quantizers with arbitrary output alphabets.

Theorem 9: Restricting to the class of (blockwise) stationary and deterministic decision rules, then there exists an asymptotically optimal quantizer ϕ that is a likelihood-ratio threshold rule.

We present the full proof of this theorem in the Appendix. The proof exploits both Lemmas 6 and 7.

VI. DISCUSSION

In this correspondence, we have studied the problem of sequential decentralized detection. For quantizers with neither local memory nor feedback (Case A in the taxonomy of Veeravalli *et al.* [22]), we have established that stationary designs need not be optimal in general. Moreover, we have shown that in the asymptotic setting (i.e., when the cost per sample goes to zero), there is a class of problems for which there exists a range of prior probabilities over which stationary strategies are suboptimal.

There are a number of open questions raised by the analysis in this correspondence. First, our analysis has established only that the best stationary rule chosen from a finite set of deterministic quantizers need not be optimal. Is there a corresponding example with an infinite number of deterministic stationary quantizer designs for which none is optimal? Second, Corollary 8 establishes the optimality of likelihood-ratio rules for randomized decision rules that produce binary outputs. This proof was based on the quasi-concavity of the function G_ϕ that specifies the asymptotic sequential cost coefficient. Is this function G_ϕ also quasi-concave for quantizers other than binary ones? Such quasi-concavity would extend the validity of Theorem 9 for the general class of randomized quantizers.

APPENDIX

Proof of Lemma 5

By renormalizing, we can assume without loss of generality (w.l.o.g.) that $a_1 + b_1 + c_1 = a_0 + b_0 + c_0 = 1$. Also w.l.o.g., assume that $b_1 \geq b_0$. Thus, $c_1 > c_0$ and $a_1 < a_0$. Replacing $c_1 = 1 - a_1 - b_1$ and $c_0 = 1 - a_0 - b_0$, the inequality $c_1/c_0 > b_1/b_0$ is equivalent to $a_1 < a_0 b_1/b_0 - (b_1 - b_0)/b_0$.

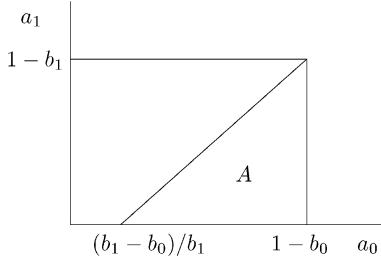
We fix values of b , and consider varying $a \in A$, where A denotes the domain for (a_0, a_1) governed by the following equality and inequality constraints: $0 < a_1 < 1 - b_1$; $0 < a_0 < 1 - b_0$; $a_1 < a_0$ and

$$a_1 < a_0 b_1/b_0 - (b_1 - b_0)/b_0. \quad (31)$$

Note that the third constraint ($a_1 < a_0$) is redundant due to the other three constraints. In particular, constraint (31) corresponds to a line passing through $((b_1 - b_0)/b_1, 0)$ and $(1 - b_0, 1 - b_1)$ in the (a_0, a_1) coordinates. As a result, A is the interior of the triangle defined by this line and two other lines given by $a_1 = 0$ and $a_0 = 1 - b_0$ (see Fig. 1).

Since both $\tilde{D}^0(a, 1-a)$ and $\tilde{D}^1(a, 1-a)$ correspond to KL divergences, they are convex functions with respect to (a_0, a_1) . In addition, the derivatives with respect to a_1 are $\frac{a_1 - a_0}{a_1(1-a_1)} < 0$ and $\log \frac{a_1(1-a_0)}{a_0(1-a_1)} < 0$, respectively. Hence, both functions can be (strictly) bounded from below by increasing a_1 while keeping a_0 unchanged, i.e., by replacing a_1 by a'_1 so that (a_0, a'_1) lies on the line given by (31), which is equivalent to the constraint $c_1/c_0 = b_1/b_0$. Let $c'_1 = 1 - b_1 - a'_1$; then $c'_1/c_0 = b_1/b_0$. Our argument has established inequalities (a) and (b) in the following chain of inequalities:

$$\tilde{D}^1(a, b+c) \quad (32a)$$

Fig. 1. Illustration of the domain A .

$$\stackrel{(a)}{>} a'_1 \log \frac{a'_1}{a_0} + (b_1 + c'_1) \log \frac{b_1 + c'_1}{b_0 + c_0} \quad (32b)$$

$$\stackrel{(b)}{=} a'_1 \log \frac{a'_1}{a_0} + c'_1 \log \frac{c'_1}{c_0} + b_1 \log \frac{b_1}{b_0} \quad (32c)$$

$$\stackrel{(c)}{\geq} (a'_1 + c'_1) \log \frac{a'_1 + c'_1}{a_0 + c_0} + b_1 \log \frac{b_1}{b_0} \quad (32d)$$

$$= \tilde{D}^1(a + c, b) \quad (32e)$$

and inequality (c) follows from an application of the log-sum inequality [9]. A similar conclusion holds for $\tilde{D}^0(a, b + c)$.

Proof of Lemma 6

Suppose the opposite is true, that there exist two sets S_1, S_2 with positive \mathbb{P}_0 -measure such that $\phi(X) = u_2$ for any $X \in S_1 \cup S_2$, and

$$\frac{f^1(S_1)}{f^0(S_1)} < \frac{f^1(u_1)}{f^0(u_1)} < \frac{f^1(S_2)}{f^0(S_2)}. \quad (33)$$

By reassigning S_1 or S_2 to the quantile u_1 , we are guaranteed to have a new quantizer ϕ' such that $D_{\phi'}^0 > D_{\phi^*}^0$ and $D_{\phi'}^1 > D_{\phi^*}^1$, thanks to Lemma 5. As a result, ϕ' has a smaller sequential cost $J_{\phi'}^*$, which is a contradiction.

Proof of Lemma 7

The proof of this lemma is conceptually straightforward, but the algebra is involved. To simplify the notation, we replace a_0 by x , a_1 by y , the function $D(a_0, a_1)$ by $f(x, y)$, and the function $D(a_1, a_0)$ by $g(x, y)$. Finally, we assume that $d_0 = d_1 = 0$; the proof will reveal that this case is sufficient to establish the more general result with arbitrary nonnegative scalars d_0 and d_1 . We have

$$f(x, y) = x \log(x/y) + (1 - x) \log[(1 - x)/(1 - y)]$$

and

$$g(x, y) = y \log(y/x) + (1 - y) \log[(1 - y)/(1 - x)].$$

Note that both f and g are convex functions and are nonnegative in their domains, and moreover that we have $F(x, y) = c_0/f(x, y) + c_1/g(x, y)$. In order to establish the quasi-concavity of F , it suffices to show that for any (x, y) in the domain of F , for any vector $h = [h_0 \ h_1] \in \mathbb{R}^2$ such that $h^T \nabla F(x, y) = 0$, there holds

$$h^T \nabla^2 F(x, y) h < 0 \quad (34)$$

(see Boyd and Vandenberghe [4]). Here we adopt the standard notation of ∇F for the gradient vector of F , and $\nabla^2 F$ for its Hessian matrix. We also use F_x to denote the partial derivative with respect to variable x , F_{xy} to denote the partial derivative with respect to x and y , and so on.

We have $\nabla F = -\frac{c_0 \nabla f}{f^2} - \frac{c_1 \nabla g}{g^2}$. Thus, it suffices to prove relation (34) for vectors of the form

$$h = \left[\left(-\frac{c_0 f_y}{f^2} - \frac{c_1 g_y}{g^2} \right) \left(\frac{c_0 f_x}{f^2} + \frac{c_1 g_x}{g^2} \right)^T \right].$$

It is convenient to write $h = c_0 v_0 + c_1 v_1$, where $v_0 = [-f_y/f^2 \ f_x/f^2]^T$ and $v_1 = [-g_y/g^2 \ g_x/g^2]^T$.

The Hessian matrix $\nabla^2 F$ can be written as $\nabla^2 F = c_0 H_0 + c_1 H_1$, where

$$H_0 = -\frac{1}{f^3} \begin{bmatrix} f_{xx} f - 2f_x^2 & f_{xy} f - 2f_x f_y \\ f_{xy} f - 2f_x f_y & f_{yy} f - 2f_y^2 \end{bmatrix}$$

and

$$H_1 = -\frac{1}{g^3} \begin{bmatrix} g_{xx} g - 2g_x^2 & g_{xy} g - 2g_x g_y \\ g_{xy} g - 2g_x g_y & g_{yy} g - 2g_y^2 \end{bmatrix}.$$

Now observe that

$$h^T \nabla^2 F h = (c_0 v_0 + c_1 v_1)^T (c_0 H_0 + c_1 H_1) (c_0 v_0 + c_1 v_1)$$

which can be simplified to

$$\begin{aligned} h^T \nabla^2 F h &= c_0^3 v_0^T H_0 v_0 + c_1^3 v_1^T H_1 v_1 \\ &\quad + c_0^2 c_1 \left(2v_0^T H_0 v_1 + v_0^T H_1 v_0 \right) + c_0 c_1^2 \left(2v_0^T H_1 v_1 + v_1^T H_0 v_0 \right). \end{aligned}$$

This function is a polynomial in c_0 and c_1 , which are restricted to be nonnegative scalars (at least one of which is assumed to be nonzero). Therefore, it suffices to prove that all the coefficients of this polynomial (with respect to c_0 and c_1) are strictly negative. In particular, we shall show that

- (i) $v_0^T H_0 v_0 \leq 0$ and
- (ii) $2v_0^T H_0 v_1 + v_0^T H_1 v_0 \leq 0$

where in both cases equality occurs only if $x = y$, which is outside of the domain of F . The strict negativity of the other two coefficients follows from entirely analogous arguments.

First, some straightforward algebra shows that inequality (i) is equivalent to the relation

$$f_{xx} f_y^2 + f_{yy} f_x^2 \geq 2f_x f_y f_{xy}.$$

But note that f is a convex function, so $f_{xx} f_{yy} \geq f_{xy}^2$. Hence, we have

$$f_{xx} f_y^2 + f_{yy} f_x^2 \stackrel{(a)}{\geq} 2\sqrt{f_{xx} f_{yy}} |f_x f_y| \stackrel{(b)}{\geq} 2f_x f_y f_{xy}$$

thereby proving (i). (In this argument, inequality (a) follows from the fact that $a^2 + b^2 \geq 2ab$, whereas inequality (b) follows from the strict convexity of f . Equality occurs only if $x = y$.) Regarding (ii), some further algebra reduces it to the inequality

$$G_1 + G_2 - G_3 \geq 0 \quad (35)$$

where

$$G_1 = 2(f_y g_y f_{xx} + f_x g_x f_{yy} - (f_y g_x + f_x g_y) f_{xy})$$

$$G_2 = f_y^2 g_{xx} + f_x^2 g_{yy} - 2f_x f_y g_{xy}$$

$$G_3 = \frac{2}{g} (f_y g_x - f_x g_y)^2.$$

At this point in the proof, we need to exploit specific information about the functions f and g , which are defined in terms of KL divergences. To simplify notation, we let $u = x/y$ and $v = (1 - x)/(1 - y)$. Computing derivatives, we have

$$\begin{aligned} f_x(x, y) &= \log(x/y) - \log((1 - x)/(1 - y)) = \log(u/v) \\ f_y(x, y) &= (1 - x)/(1 - y) - x/y = v - u \\ g_x(x, y) &= (1 - y)/(1 - x) - y/x = 1/v - 1/u \\ g_y(x, y) &= \log(y/x) - \log((1 - y)/(1 - x)) = \log(v/u) \\ \nabla^2 f(x, y) &= \begin{bmatrix} \frac{1}{x(1-x)} & -\frac{1}{y(1-y)} \\ -\frac{1}{x(1-x)} & \frac{1}{(1-y)^2} + \frac{1}{y^2} \end{bmatrix} \\ \nabla^2 g(x, y) &= \begin{bmatrix} \frac{1}{(1-x)^2} + \frac{1}{x^2} & -\frac{1}{x(1-x)} \\ -\frac{1}{x(1-x)} & \frac{1}{y(1-y)} \end{bmatrix}. \end{aligned}$$

Noting that $f_x = -g_y$; $g_{xy} = -f_{xx}$; $f_{xy} = -g_{yy}$, we see that (35) is equivalent to

$$2(f_x g_x f_{yy} + f_y g_x g_{yy}) - f_x^2 g_{yy} + f_y^2 g_{xx} \geq \frac{2}{g} (f_y g_x - f_x g_y)^2. \quad (36)$$

To simplify the algebra further, we shall make use of the inequality $(\log t^2)^2 \leq (t - 1/t)^2$, which is valid for any t . This implies that

$$\begin{aligned} f_y g_x &= (v - u)(1/v - 1/u) \leq f_x g_y = -(\log(u/v))^2 \\ &= -f_x^2 = -g_y^2 \leq 0. \end{aligned}$$

Thus, $-f_x^2 g_{yy} \geq f_y g_x g_{yy}$ and

$$\frac{2}{g} (f_y g_x - f_x g_y)^2 \leq \frac{2}{g} f_y g_x (f_y g_x - f_x g_y).$$

As a result, (36) would follow if we can show that

$$2(f_x g_x f_{yy} + f_y g_x g_{yy}) + f_y g_x g_{yy} + f_y^2 g_{xx} \geq \frac{2}{g} f_y g_x (f_y g_x - f_x g_y).$$

For all $x \neq y$, we may divide both sides by $-f_y(x, y)g_x(x, y) > 0$. Consequently, it suffices to show that

$$-2f_x f_{yy}/f_y - f_y g_{xx}/g_x - 3g_{yy} \geq \frac{2}{g} (f_x g_y - g_x f_y)$$

or, equivalently

$$\begin{aligned} 2 \log(u/v) \left(\frac{v}{u-1} + \frac{u}{1-v} \right) + \left(\frac{u}{1-x} + \frac{v}{x} \right) \\ - \frac{3}{y(1-y)} \geq \frac{2}{g} \left(\frac{(u-v)^2}{uv} - \left(\log \frac{u}{v} \right)^2 \right) \end{aligned}$$

or, equivalently

$$\begin{aligned} 2 \log(u/v) \frac{(u-v)(u+v-1)}{(u-1)(1-v)} + \frac{(u-v)^2(u+v-4uv)}{uv(u-1)(1-v)} \\ \geq \frac{2}{g} \left(\frac{(u-v)^2}{uv} - \left(\log \frac{u}{v} \right)^2 \right). \quad (37) \end{aligned}$$

Due to the symmetry, it suffices to prove (37) for $x < y$. In particular, we shall use the following inequality for logarithm mean [13], which holds for $u \neq v$:

$$\frac{3}{2\sqrt{uv} + (u+v)/2} < \frac{\log u - \log v}{u-v} < \frac{1}{(uv(u+v)/2)^{1/3}}.$$

We shall replace $\frac{\log(u/v)}{u-v}$ in (37) by appropriate upper and lower bounds. In addition, we shall also bound $g(x, y)$ from below, using the following argument. When $x < y$, we have $u < 1 < v$, and

$$\begin{aligned} g(x, y) &= y \log \frac{y}{x} + (1-y) \log \frac{1-y}{1-x} \\ &> \frac{3y(y-x)}{2\sqrt{xy} + (x+y)/2} + \frac{(1-y)(x-y)}{[(1-x)(1-y)(1-(x+y)/2)]^{1/3}} \\ &= \frac{3(1-v)(1-u)}{(u-v)(2\sqrt{u} + \frac{u+1}{2})} + \frac{(u-1)(1-v)}{(u-v)(v(v+1)/2)^{1/3}} > 0. \end{aligned}$$

Let us denote this lower bound by $g(u, v)$.

Having got rid of the logarithm terms, (37) will hold if we can prove the following:

$$\begin{aligned} \frac{6(u-v)^2(u+v-1)}{(2\sqrt{uv} + (u+v)/2)(u-1)(1-v)} + \frac{(u-v)^2(u+v-4uv)}{uv(u-1)(1-v)} \\ \geq \frac{2}{g(u, v)} \left(\frac{(u-v)^2}{uv} - \frac{9(u-v)^2}{(2\sqrt{uv} + (u+v)/2)^2} \right) \end{aligned}$$

or equivalently,

$$\begin{aligned} \left(\frac{6(u+v-1)}{(2\sqrt{uv} + (u+v)/2)} + \frac{(u+v-4uv)}{uv} \right) \\ \times \left(\frac{3}{(v-u)(2\sqrt{u} + \frac{u+1}{2})} - \frac{1}{(v-u)(v(v+1)/2)^{1/3}} \right) \\ \geq 2 \left(\frac{1}{uv} - \frac{9}{(2\sqrt{uv} + (u+v)/2)^2} \right) \quad (38) \end{aligned}$$

which is equivalent to

$$\begin{aligned} \frac{(u+v-2\sqrt{uv})((u+v)/2 + 3\sqrt{uv} + 4uv)}{(2\sqrt{uv} + (u+v)/2)uv} \\ \times \frac{3(v(v+1)/2)^{1/3} - (2\sqrt{u} + (u+1)/2)}{(v-u)(2\sqrt{u} + (u+1)/2)(v(v+1)/2)^{1/3}} \\ \geq \frac{(u+v-2\sqrt{uv})((u+v)/2 + 5\sqrt{uv})}{uv(2\sqrt{uv} + (u+v)/2)^2} \quad (39) \end{aligned}$$

and also equivalent to

$$\begin{aligned} ((u+v)/2 + 2\sqrt{uv})((u+v)/2 + 3\sqrt{uv} + 4uv) \\ \times [3(v(v+1)/2)^{1/3} - (2\sqrt{u} + (u+1)/2)] \\ \geq (2\sqrt{u} + (u+1)/2)(v(v+1)/2)^{1/3} \\ \times ((u+v)/2 + 5\sqrt{uv})(v-u). \quad (40) \end{aligned}$$

It can be checked by tedious but straightforward calculus that inequality (40) holds for any $u \leq 1 \leq v$, and equality holds when $u = 1 = v$, i.e., $x = y$.

Proof of Theorem 9

Suppose that ϕ is not a likelihood-ratio rule. Then there exist positive \mathbb{P}_0 -probability disjoint sets S_1, S_2, S_3 such that for any $X_1 \in S_1, X_2 \in S_2, X_3 \in S_3$

$$\phi(X_1) = \phi(X_3) = u_1 \quad (41a)$$

$$\phi(X_2) = u_2 \neq u_1 \quad (41b)$$

$$\frac{f^1(X_1)}{f^0(X_1)} < \frac{f^1(X_2)}{f^0(X_2)} < \frac{f^1(X_3)}{f^0(X_3)}. \quad (41c)$$

Define the probability of the quantiles as

$$f^0(u_1) := \mathbb{P}_0(\phi(X) = u_1),$$

$$f^0(u_2) := \mathbb{P}_0(\phi(X) = u_2);$$

$$f^1(u_1) := \mathbb{P}_1(\phi(X) = u_1),$$

$$f^1(u_2) := \mathbb{P}_1(\phi(X) = u_2).$$

Similarly, for the sets S_1, S_2 , and S_3 , we define

$$a_0 = f^0(S_1), \quad b_0 = f^0(S_2), \quad \text{and} \quad c_0 = f^0(S_3)$$

$$a_1 = f^1(S_1), \quad b_1 = f^1(S_2), \quad \text{and} \quad c_1 = f^1(S_3).$$

Finally, let p_0, p_1, q_0 , and q_1 denote the probability measures of the ‘‘residuals’’

$$p_0 = f^0(u_2) - b_0, \quad p_1 = f^1(u_2) - b_1,$$

$$q_0 = f^0(u_1) - a_0 - c_0, \quad q_1 = f^1(u_1) - a_1 - c_1.$$

Note that we have $\frac{a_1}{a_0} < \frac{b_1}{b_0} < \frac{c_1}{c_0}$. In addition, the sets S_1 and S_3 were chosen so that $\frac{a_1}{a_0} < \frac{a_1}{a_0} < \frac{a_1}{a_0} < \frac{c_1}{c_0}$. From Lemma 6, there holds

$$\frac{p_1 + b_1}{p_0 + b_0} = \frac{f^1(u_2)}{f^0(u_2)} \notin \left(\frac{a_1}{a_0}, \frac{c_1}{c_0} \right).$$

We may assume w.l.o.g. that $\frac{p_1+b_1}{p_0+b_0} \leq \frac{a_1}{a_0}$. Then, $\frac{p_1+b_1}{p_0+b_0} < \frac{b_1}{b_0}$, so $\frac{p_1}{p_0} < \frac{p_1+b_1}{p_0+b_0}$. Overall, we are guaranteed to have the ordering

$$\frac{p_1}{p_0} < \frac{p_1+b_1}{p_0+b_0} \leq \frac{a_1}{a_0} < \frac{b_1}{b_0} < \frac{c_1}{c_0}. \quad (42)$$

Our strategy will be to modify the quantizer ϕ only for those X for which $\phi(X)$ takes the values u_1 or u_2 , such that the resulting quantizer is defined by a LLR-based threshold, and has a smaller (or equal) value of the corresponding cost J_ϕ^* . For simplicity in notation, we use \mathcal{A} to denote the set with measures under \mathbb{P}_0 and \mathbb{P}_1 equal to a_0 and a_1 ; the sets \mathcal{B} , \mathcal{C} , \mathcal{P} , and \mathcal{Q} are defined in an analogous manner. We begin by observing that we have either $\frac{a_1}{a_0} \leq \frac{q_1+a_1}{q_0+a_0} < \frac{b_1}{b_0}$ or $\frac{b_1}{b_0} < \frac{q_1+c_1}{q_0+c_0} \leq \frac{c_1}{c_0}$. Thus, in our subsequent manipulation of sets, we always bundle \mathcal{Q} with either \mathcal{A} or \mathcal{C} accordingly without changing the ordering of the probability ratios. Without loss of generality, then, we may disregard the corresponding residual set corresponding to \mathcal{Q} in the analysis to follow.

In the remainder of the proof, we shall show that either one of the following two modifications of the quantizer ϕ will improve (decrease) the sequential cost J_ϕ^* .

- (i) Assign \mathcal{A} , \mathcal{B} and \mathcal{C} to the same quantization level u_1 , and leave \mathcal{P} to the level u_2 ; or
- (ii) Assign \mathcal{P} , \mathcal{A} and \mathcal{B} to the same level u_2 , and leave \mathcal{C} to the level u_1 .

It is clear that this modified quantizer design respects the likelihood-ratio rule for the quantization indices u_1 and u_2 . By repeated application of this modification for every such pair, we are guaranteed to arrive at a likelihood-ratio quantizer that is optimal, thereby completing the proof.

Let a'_0, b'_0, c'_0, p'_0 be normalized versions of a_0, b_0, c_0, p_0 , respectively (i.e., $a'_0 = a_0/(p_0 + a_0 + b_0 + c_0)$, and so on). Similarly, let a'_1, b'_1, c'_1, p'_1 be normalized versions of a_1, b_1, c_1, p_1 , respectively. With this notation, we have the relations

$$\begin{aligned} D_\phi^0 &= \sum_{u \neq u_1, u_2} f^0(u) \log \frac{f^0(u)}{f^1(u)} \\ &\quad + (p_0 + b_0) \log \frac{p_0 + b_0}{p_1 + b_1} + (a_0 + c_0) \log \frac{a_0 + c_0}{a_1 + c_1} \\ &= A_0 + (f^0(u_1) + f^0(u_2)) \\ &\quad \times \left((p'_0 + b'_0) \log \frac{p'_0 + b'_0}{p'_1 + b'_1} + (a'_0 + c'_0) \log \frac{a'_0 + c'_0}{a'_1 + c'_1} \right) \\ &= A_0 + (f^0(u_1) + f^0(u_2)) \tilde{D}^0(p' + b', a' + c'), \\ D_\phi^1 &= \sum_{u \neq u_1, u_2} f^1(u) \log \frac{f^1(u)}{f^0(u)} \\ &\quad + (p_1 + b_1) \log \frac{p_1 + b_1}{p_0 + b_0} + (a_1 + c_1) \log \frac{a_1 + c_1}{a_0 + c_0} \\ &= A_1 + (f^1(u_1) + f^1(u_2)) \tilde{D}^1(p' + b', a' + c') \end{aligned}$$

where we define

$$\begin{aligned} A_0 &:= \sum_{u \neq u_1, u_2} f^0(u) \log \frac{f^0(u)}{f^1(u)} \\ &\quad + (f^0(u_1) + f^0(u_2)) \log \frac{f^0(u_1) + f^0(u_2)}{f^1(u_1) + f^1(u_2)} \geq 0 \\ A_1 &:= \sum_{u \neq u_1, u_2} f^1(u) \log \frac{f^1(u)}{f^0(u)} \\ &\quad + (f^1(u_1) + f^1(u_2)) \log \frac{f^1(u_1) + f^1(u_2)}{f^0(u_1) + f^0(u_2)} \geq 0 \end{aligned}$$

due to the nonnegativity of the KL divergences.

Note that from (42) we have

$$\frac{p'_1}{p'_0} < \frac{p'_1 + b'_1}{p'_0 + b'_0} \leq \frac{a'_1}{a'_0} < \frac{b'_1}{b'_0} < \frac{c'_1}{c'_0}$$

in addition to the normalization constraints that $p'_0 + a'_0 + b'_0 + c'_0 = p'_1 + a'_1 + b'_1 + c'_1 = 1$. It follows that

$$\frac{p'_1 + b'_1}{p'_0 + b'_0} < \frac{p'_1 + a'_1 + b'_1 + c'_1}{p'_0 + a'_0 + b'_0 + c'_0} = 1.$$

Let us consider varying the values of a'_1, b'_1 , while fixing all other variables and ensuring that all the above constraints hold. Then, $a'_1 + b'_1$ is constant, and both $\tilde{D}^0(p' + b', a' + c')$ and $\tilde{D}^1(p' + b', a' + c')$ increase as b'_1 decreases and a'_1 increases. In other words, if we define $a''_0 = a'_0, b''_0 = b'_0$ and a''_1 and b''_1 such that

$$\frac{a''_1}{a''_0} = \frac{b''_1}{b''_0} = \frac{1 - p'_1 - c'_1}{1 - p'_0 - c'_0}$$

then we have

$$\begin{aligned} \tilde{D}^0(p' + b', a' + c') &\leq \tilde{D}^0(p' + b'', a'' + c') \text{ and} \\ \tilde{D}^1(p' + b', a' + c') &\leq \tilde{D}^1(p' + b'', a'' + c'). \end{aligned} \quad (43)$$

Now note that vector (b''_0, b''_1) in \mathbb{R}^2 is a convex combination of $(0, 0)$ and $(a''_0 + b''_0, a''_1 + b''_1)$. It follows that $(p'_0 + b''_0, p'_1 + b''_1)$ is a convex combination of (p'_0, p'_1) and $(p'_0 + a''_0 + b''_0, p'_1 + a''_1 + b''_1) = (p'_0 + a'_0 + b'_0, p'_1 + a'_1 + b'_1)$. By (43), we obtain

$$\begin{aligned} G_\phi &= \frac{\pi^0}{A_0 + (f^0(u_1) + f^0(u_2)) \tilde{D}^0(p' + b', a' + c')} \\ &\quad + \frac{\pi^1}{A_1 + (f^1(u_1) + f^1(u_2)) \tilde{D}^1(p' + b', a' + c')} \\ &\geq \frac{\pi^0}{A_0 + (f^0(u_1) + f^0(u_2)) \tilde{D}^0(p' + b'', a'' + c')} \\ &\quad + \frac{\pi^1}{A_1 + (f^1(u_1) + f^1(u_2)) \tilde{D}^1(p' + b'', a'' + c')} \\ &= \frac{\pi^0}{A_0 + (f^0(u_1) + f^0(u_2)) D(p'_0 + b'_0, p'_1 + b'_1)} \\ &\quad + \frac{\pi^1}{A_1 + (f^1(u_1) + f^1(u_2)) D(p'_1 + b'_1, p'_0 + b'_0)}. \end{aligned}$$

Applying the quasi-concavity result in Lemma 7

$$\begin{aligned} G_\phi &\geq \min \left\{ \frac{\pi^0}{A_0 + (f^0(u_1) + f^0(u_2)) D(p'_0, p'_1)} \right. \\ &\quad + \frac{\pi^1}{A_1 + (f^1(u_1) + f^1(u_2)) D(p'_1, p'_0)}, \\ &\quad \frac{\pi^0}{A_0 + (f^0(u_1) + f^0(u_2)) D(p'_0 + a'_0 + b'_0, p'_1 + a'_1 + b'_1)} \\ &\quad \left. + \frac{\pi^1}{A_1 + (f^1(u_1) + f^1(u_2)) D(p'_1 + a'_1 + b'_1, p'_0 + a'_0 + b'_0)} \right\}. \end{aligned}$$

But the two arguments of the minimum are the sequential cost coefficient corresponding to the two possible modifications of ϕ . Hence, the proof is complete.

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Myopic Coding in Multiterminal Networks

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Abstract—This correspondence investigates the interplay between cooperation and achievable rates in multiterminal networks. Cooperation refers to the process of nodes working together to relay data toward the destination. There is an inherent tradeoff between achievable information transmission rates and the level of cooperation, which is determined by how many nodes are involved and how the nodes encode/decode the data. We illustrate this tradeoff by studying information-theoretic decode-forward-based coding strategies for data transmission in multiterminal networks. Decode-forward strategies are usually discussed in the context of omniscient coding, in which all nodes in the network fully cooperate with each other, both in encoding and decoding. In this correspondence, we investigate myopic coding, in which each node cooperates with only a few neighboring nodes. We show that achievable rates of myopic decode-forward can be as large as that of omniscient decode-forward in the low signal-to-noise ratio (SNR) regime. We also show that when each node has only a few cooperating neighbors, adding one node into the cooperation increases the transmission rate significantly. Furthermore, we show that myopic decode-forward can achieve nonzero rates as the network size grows without bound.

Index Terms—Achievable rates, decode-forward, multiple-relay channel, multiterminal network, myopic coding.

I. INTRODUCTION

A. Wireless Networks

Wireless networks have been receiving much attention recently by both researchers and industry. The main advantage of wireless technology to users is the seamless access to the network whenever and wherever they are; to service providers, easier deployment, as no cable laying is required. Examples of wireless networks include cellular mobile networks, Wi-Fi networks, and sensor networks. A large amount of research has been carried out recently on various aspects of wireless networks, including power saving [1], [2], routing [3]–[5], transport capacity [6], [7], and connectivity [8]. In this correspondence, we focus on transmission rates in multiterminal wireless networks.

Analyzing transmission rates in multiterminal networks is not easy. Consider the *single-relay channel* [9], [10], a channel consisting of one source, one relay, and one destination. Even for this simple three-terminal network, the capacity is not known except for a few special cases, e.g., the degraded relay channel [9]. This hints at the difficulty of analyzing multiterminal networks. We attempt to investigate an excerpt of the multiterminal network by looking at data transmission from a single source to a single destination, from multiple sources to a single destination, and from a single source to multiple destinations, with the help of relay(s). Appropriate models for these types of networks are the *multiple-relay channel* [11], [12] (an extension of the single-relay channel), the *multiple-access relay channel* [13], [14], and the *broadcast relay channel* [15], respectively. The reason for using relays, which have no

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